# Numbers 

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## Outline

(1) Binary numbers
(2) Hexadecimal numbers
(3) Integers
(4) Floating point numbers
(5) Converting decimal numbers to floats
(6) Floating point mathematics

## Binary numbers

- Decimal place value system

$$
\begin{aligned}
15301201 & =1 * 10^{7}+5 * 10^{6}+3 * 10^{5}+10^{3}+2 * 10^{2}+1 \\
& =10000000+5000000+300000+1000+200+1 \\
& =15301201
\end{aligned}
$$

- Binary place value system

$$
\begin{aligned}
10101111 & =2^{7}+2^{5}+2^{3}+2^{2}+2+1 \\
& =128+32+8+4+2+1 \\
& =175
\end{aligned}
$$

## Bit numbering



- The least significant bit of a byte is bit 0
- The most significant bit is bit 7
- In yasm this number could be written as 10101111b


## Decimal to binary conversion

- Convert 741 to binary
- Repeatedly divide by 2 and keep the remainders

| division |  | remainder | bits |
| ---: | :--- | :---: | ---: |
| $741 / 2$ | $=370$ | 1 | 1 |
| $370 / 2$ | $=185$ | 0 | 01 |
| $185 / 2=92$ | 1 | 101 |  |
| $92 / 2=46$ | 0 | 0101 |  |
| $46 / 2=23$ | 0 | 00101 |  |
| $23 / 2=11$ | 1 | 100101 |  |
| $11 / 2=5$ | 1 | 1100101 |  |
| $5 / 2=2$ | 1 | 11100101 |  |
| $2 / 2=1$ | 0 | 011100101 |  |
| $1 / 2=0$ | 1 | 1011100101 |  |

## Hexadecimal numbers

- Base 16 numbers
- Use as "digits" 0-9 and A-F (or a-f)
- $A=10, B=11, C=12, D=13, E=14, F=15$

$$
\begin{aligned}
\text { 0x2b1a } & =2 * 16^{3}+11 * 16^{2}+1 * 16+10 \\
& =2 * 4096+11 * 256+16+10 \\
& =8192+2816+16+10 \\
& =11034
\end{aligned}
$$

## Why use hexadecimal?

- Each hexadecimal digit or "nibble" is 4 bits
- 0x2b1a $=0010101100011010$
- $0 x 2 b 1 a=0010101100011010 b$
- Counting 32 bits for a binary pattern would be hard
- Hexadecimal is much easier
- $0 x d e a d b e e f=11011110101011011011111011101111 \mathrm{~b}$


## Converting decimal to hexadecimal

- Convert 40007 to hexadecimal
- Repeatedly divide by 16 and keep the remainders

| division | remainder | hex |  |
| ---: | :--- | :--- | ---: |
| $40007 / 16$ | $=2500$ | 7 | 7 |
| $2500 / 16$ | $=156$ | 4 | 47 |
| $156 / 16$ | $=9$ | 12 | $c 47$ |
| $9 / 16$ | $=0$ | 9 | $9 c 47$ |

## Integers

- Integers can be 1, 2, 4 or 8 bytes long
- They can be signed or unsigned

| Variety | Bits | Bytes | Minimum | Maximum |
| :---: | :---: | :---: | :---: | :---: |
| unsigned | 8 | 1 | 0 | 255 |
| signed | 8 | 1 | -128 | 127 |
| unsigned | 16 | 2 | 0 | 65535 |
| signed | 16 | 2 | -32768 | 32767 |
| unsigned | 32 | 4 | 0 | 4294967295 |
| signed | 32 | 4 | -2147483648 | 2147483647 |
| unsigned | 64 | 8 | 0 | 18446744073709551615 |
| signed | 64 | 8 | -9223372036854775808 | 9223372036854775807 |

## Negative integers

- We use the highest-order bit as a sign bit
- 1 for a sign bit means a negative number
- If we stored -1 as 10000001b
- $-1+1$ would be $10000001 b+00000001 b=100000010 b$
- Then addition would yield $-1+1=-2$
- There must be a better way to store negatives
- Hopefully, we can use the same circuitry for positives and negatives


## Two's complement integers

- To convert a number to its negative, use two's complement
- Flip all the bits
- Add 1
- Let's convert 1 to -1 with 8 bit numbers

> 00000001 for the absolute value
> 11111110 for the complement
> 11111111 after adding 1 to the complement $-1=11111111$

- Two's complement negative numbers work for addition


## More 8 bit signed integers

- They form a cycle if you keep adding 1

```
00000000 = 0
00000001 = 1
00000010 = 2
01111111 = 127
10000000 = -128
10000001 = -127
10000010 = -126
11111110 = -2
11111111 = -1
00000000 = 0
```


## Addition

- Let's convert and add -29124 + 125

```
29124 = 0111000111000100
Negate = 1000111000111011
Add 1 = 1000111000111100
    125 = 0000000001111101
Now add 1000111000111100
        0000000001111101
    1000111010111001
Negate 0111000101000110
Add 1 0111000101000111
                28999
So -29124 + 125 = -28999
```


## Binary multiplication

|  | 1010101 |
| ---: | ---: |
| $*$ | 10101 |
|  | 1010101 |
| 1010101 |  |
|  | 1010101 |
|  | 11011111001 |

## Floating point numbers

- 32 bit, 64 bit and 80 bit numbers
- Stored in IEEE 754 format

| Variety | Bits | Exponent | Exponent Bias | Fraction | Precision |
| :--- | :---: | :---: | :---: | :---: | :---: |
| float | 32 | 8 | 127 | 23 | $\sim 7$ digits |
| double | 64 | 11 | 1023 | 52 | $\sim 16$ digits |
| long double | 80 | 15 | 16383 | 64 | 19 digits |

- Exponents are binary exponents
- An exponent field has the bias added
- A 32 bit exponent field of 128 means a binary exponent 1
- A 32 bit exponent field of 125 means a binary exponent -2
- 0.0 is stored as all bits equal to 0
- Exponent field 255 means "Not a Number"


## Binary numbers with binary points

$$
\begin{aligned}
0.1_{2} & =2^{-1} \\
& =0.5 \\
1.11_{2} & =1+2^{-1}+2^{-2} \\
& =1+0.5+0.25 \\
& =1.75 \\
1001.1001_{2} & =2^{3}+1+2^{-1}+2^{-4} \\
& =8+1+0.5+0.0625 \\
& =9.5625 \\
1.0010101 * 2^{3} & =1001.0101 \\
& =2^{3}+1+2^{-2}+2^{-4} \\
& =8+1+0.25+0.0625 \\
& =9.3125
\end{aligned}
$$

## Implicit 1 bit



- Normalized floats have exponent fields from 1 to 254
- For these floats there will be at least one 1 bit in the number
- IEEE 754 uses implicit 1 bits
- For non-zero floats, they can be written in "scientific" notation
- $1011.10101=1.01110101 * 2^{3}$
- The leading 1 bit is not stored
- The value (fraction) field is 01110101000000000000000
- So we have 23 bits of fraction with 1 implicit bit $=24$ bits
- The sign bit is flipped to negate a float (1 means negative)


## Floating point storage

- Consider consider this listing by yasm

1
2
30000000000000000
$4000000040000803 F$
500000008 000080BF
6 0000000C 0000E03F
700000010 0000F542
800000014 CDCC8C3F
900000018 F9021550
\%line $1+1$ fp.asm
[section .data]
zero dd 0.0
one dd 1.0
neg1 dd -1.0
a dd 1.75
b dd 122.5
d dd 1.1
e dd 10000000000.0

- The bytes are backwards
- 1.0 should be represented logically as 3F800000
- 0 sign bit, 127 exponent field, 0 for the fraction field


## Floating point storage (2)

$4000000040000803 F$
500000008 000080BF
6 0000000C 0000E03F
700000010 0000F542

```
one dd 1.0
neg1 dd -1.0
a dd 1.75
b dd 122.5
```

- All these have a lot of 0 bits in the fractions
- They are all exactly equal to a sum of a few powers of 2
- $1=2^{0}$
- $1.75=2^{0}+2^{-1}+2^{-2}$
- $122.5=2^{6}+2^{5}+2^{4}+2^{3}+2^{1}+2^{-1}$
- -1.0 differs from 1.0 only in the sign bit


## Floating point storage (3)

800000014 CDCC8C3F d dd 1.1

- 1.1 is a repeating binary number
- The number in "proper" order is 3F8CCCCD
- The exponent field is 127 , so the exponent is 1
- The number is $1.00011001100110011001101_{2}$
- It looks like $1.1=1.000 \overline{1100}$


## Converting decimal numbers to floats

- Determine the sign bit and work with the absolute value
- Convert the whole part of the decimal number
- Convert the fraction
- Express in binary scientific notation
- Build the exponent field by adding 127 bias
- Drop the leading 1 to get the fraction field
- Example: convert -12.25
- Sign bit is 1
- Whole part is $12=1100_{2}$
- Fraction is $0.25=0.01$
- Scientific notation $12.25=1.10001_{2} * 2^{3}$

$$
\begin{aligned}
-12.25 & =11000001010001000000000000000000 \\
& =0 x C 1440000
\end{aligned}
$$

## Converting decimal number to float (2)

- The only non-obvious step is coverting the fractional part to a binary fraction.
- Suppose you have a decimal number $x=$.abcdefgh
- Then if you multiple $x$ by 2 , the only possible result is $2 x<1$ or $1 \leq 2 x<2$
- If $2 x<1$, then $x<0.5$, which means the first bit after the binary point is 0 .
- If $2 x \geq 1$, then $x \geq 0.5$, which means the first bit after the binary point is 1 .
- So we set the first bit and work on the remaining fractional part of $2 x$ to get the next bit.
- This process continues until we reach $x=0$ or we have enough bits.


## Converting decimal number to float (3)

- Let's convert -121.6875 to a binary number
- First the sign is 0
- $121=1111001_{2}$
- Now it's time to work on 6875

| Multiply |  | Result | Binary |
| :--- | :--- | :--- | :--- |
| $.6875 * 2$ | $=$ | 1.375 | $.1_{2}$ |
| $.375 * 2$ | $=$ | 0.75 | $.10_{2}$ |
| $.75 * 2$ | $=$ | 1.5 | $.101_{2}$ |
| $.5 * 2$ | $=1.0$ | $.1011_{2}$ |  |

- $-121.6875=-1111001.1011_{2}$
- $-121.6875=-1.1110011011_{2} * 2^{6}$
- As a binary float 11000010111100110110000000000000
- Expressed in hexadecimal: 0xC2F36000


## Floating point addition

- Let's add 41.275 and 0.315
- $41.275=101001.010001100110011010$ in binary
- $0.325=0.0101000010100011110101110$ in binary
- As with decimals, we align the numbers and add

| 101001.010001100110011010 |
| ---: |
| $+\quad 0.0101000010100011110101110$ |
| 101001.1001011100001010010101110 |

- There are 31 digits in the answer
- The answer must be rounded to 24 bits
- Rounding the last 7 bits means truncation in this case
- We get $0 \times 42265 \mathrm{c} 29$ which is 41.59 (approximately)


## Floating point multiplication

- Let's multiply 7.5 and 4.375

|  | 7.5 | $=$ |
| ---: | ---: | ---: |
| $*$ | $11111_{2}$ |  |
| 4.375 | $=$ | $100.011_{2}$ |
|  |  |  |
|  |  | $11111_{2}$ |
|  |  | $111100000_{2}$ |
|  |  |  |
|  |  | $100000.1101_{2}$ |

- Conversion to float format should be apparent by now

